

# Stability Region of Wireless Networks

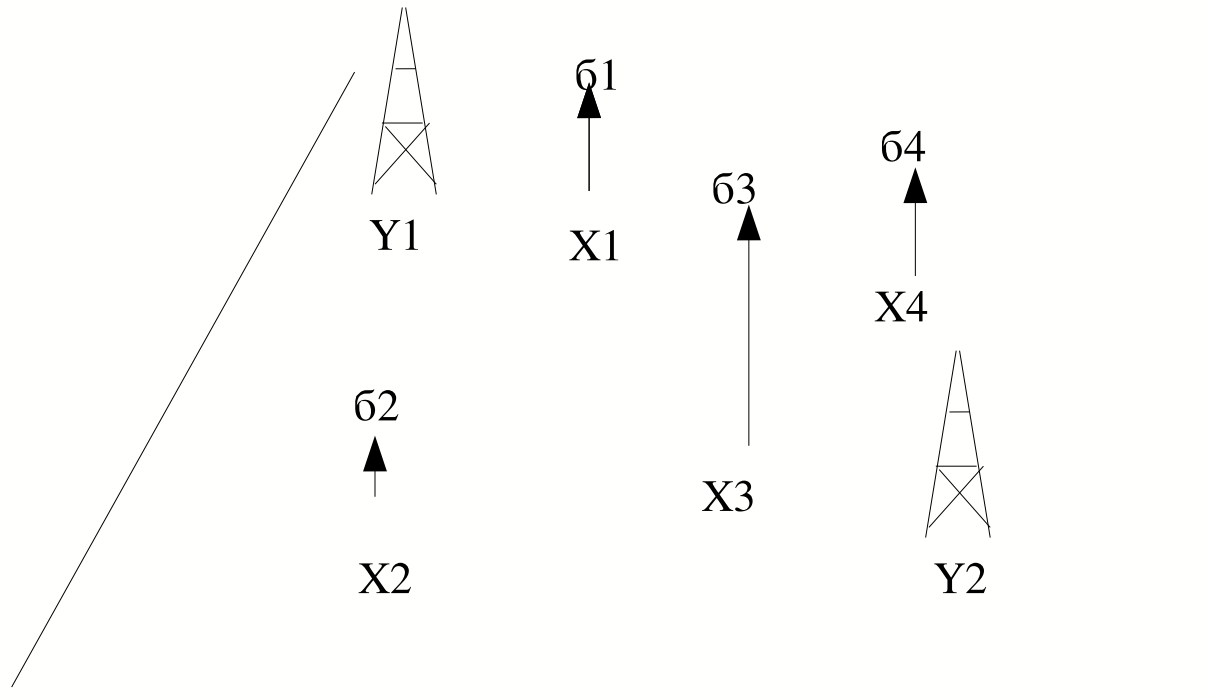
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## DATA FLOWS ON THE DOWNLINK

- Users arrive in a wireless network and leave when they have received some data.
- Processing rates vary with positions.
- We consider a large network deployed in the plane.



## MODEL DESCRIPTION : ARRIVAL POINT PROCESS

- The  $n^{\text{th}}$  user arrives at time  $T_n$  located at  $X_n \in \mathbb{R}^2$  and requires a service of  $\sigma_n$  bits.
- Customers marked point process  $\mathbb{R} \times \mathbb{R}^2 \times \mathbb{R}_+$  :

$$A = \sum_n \delta_{\{T_n, X_n, \sigma_n\}},$$

- $A$  is **time-ergodic** (but not necessarily space-stationary)

## MODEL DESCRIPTION : ARRIVAL POINT PROCESS

- The intensity measure of  $A$ ,  $\lambda(dx)dt$ , is a **Radon measure** and

$$\sigma(x) = E^{0,x}(\sigma_{0,x}) \in \mathbb{R}_+.$$

$\sigma(x)$  = "mean bits requirement of a typical user located at  $x$ "

For a set  $B$ ,  $\lambda(B)$  = "mean number of users arriving by time unit in  $B$ "

- For  $t \geq T_n$ , let  $\sigma_n(t)$  be the remaining service of the  $n^{\text{th}}$  user. The **workload at time  $t$**  is the atomic measure :

$$W_t = \sum_{n:T_n \leq t} \sigma_n(t) \delta_{X_n}.$$

## MODEL DESCRIPTION : SERVER STATIONS

- A countable set of server stations with a processing rate  $x \rightarrow r_j(x)$  such that :

$$\lim_{|x| \rightarrow +\infty} r_j(x) = 0.$$

- The server stations are in a **random environment**. The processing power available at time  $t$  is :

$$P_j(t) \geq 0, \quad P_j \text{ is ergodic} \quad \text{and} \quad p_j = EP_j < \infty$$

- We suppose : for all bounded set  $B$ ,  $\sum_j p_j \text{ess sup}_{x \in B} r_j(x) < +\infty$ .

## MODEL DESCRIPTION : ADAPTIVE POLICY

**Definition 1.** Let  $\mathcal{M}$  be the set of Radon measures on  $\mathbb{R}^2$ , the *policy enforced by  $j^{\text{th}}$  the server station* is a mapping :

$$\begin{aligned}\mathcal{M} &\rightarrow \mathcal{M} \\ m &\mapsto \pi_j(m).\end{aligned}$$

such that :

$$\pi_j(\mathbb{R}^2) = \int_{\mathbb{R}^2} \pi_j(m)(dx) \leq 1$$

and

$\pi_j(m)$  is absolutely continuous with respect to  $m$

The policy enforced at time  $t$  is :  $\pi_j(t) = \pi_j(W_t)$  .

## MODEL DESCRIPTION : EVOLUTION EQUATION

We can define the dynamic of our queueing system, for all sets  $B$  and  $t' > t$  :

$$W_{t'}(B) = W_t(B) + \int_t^{t'} \int_B \sigma_{s,x} A(ds \times dx) - \sum_j \int_t^{t'} \int_B P_j(s) r_j(x) \pi_j(s)(dx) ds,$$

where :  $\int_t^{t'} \int_B \sigma_{s,x} A(ds \times dx) = \sum_n \sigma_n \mathbb{I}(T_n \in [t, t'[) \mathbb{I}(X_n \in B)$ .

$$W_{t'} = W_t + \Sigma(t, t') - \sum_j \int_t^{t'} P_j(s) r_j \pi_j(s) ds.$$

## TDMA CELLULAR NETWORK

- Server station  $j$  is located at  $Y_j \in \mathbb{R}^2$ .
- To server station  $j$  is attached a cell  $V_j$  with  $V_j \cap V_{j'} = \emptyset$  if  $j \neq j'$  and  $\cup_j V_j = \mathbb{R}^2$ . Such collection of sets  $\{V_j\}_j$  will be called a **tessellation**.
- A server station serves the users in its cell :

$$\forall m \in \mathcal{M}, \quad \pi_j(m)(\mathbb{R}^2 \setminus V_j) = 0.$$

- $P_j(t) = 1$  and a server station serves one user at a time :

$$\pi_j(m) = \delta_{x_j} \text{ or } 0.$$

## TDMA CELLULAR NETWORK

- The server stations are constantly emitting at a power  $S$
- Communication channel from  $x$  to  $y$  with attenuation function  $L(x, y)$ , (for example  $L(x, y) = |x - y|^{-\alpha}$ ,  $\alpha > 2$ ).
- Shot Noise :  $I(x) = \sum_k L(x, Y_k)$ .
- Gaussian white noise at reception of spectral density  $\eta$ ,
- Available bandwidth of  $\Delta$  Hz.

$$r_j(x) = \Delta \log_2 \left( 1 + \frac{SL(x, Y_j)}{\Delta\eta + S(I(x) - L(x, Y_j))} \right).$$

## TDMA CELLULAR NETWORK : FAST FADING

- $r_j$  changes over time. We suppose that  $r_j \in \{r_{j,1}, \dots, r_{j,k}, \dots\}$ .
- We divide server station  $j$  into servers  $\{(j, k)\}_{k \in \mathbb{N}}$ .
- $P_{j,k}(t) \in \{0, 1\}$ , the server stations is switched on or switched off and :

$$\sum_k P_{j,k}(t) = 1.$$

One server station is switched on at any time.

## CDMA NETWORK IN MACRODIVERSITY

- The network is in **macrodiversity** : the server stations may cooperate to send data to users.
- At time  $t$ , the signal emitted by the  $j^{th}$  base station for the  $n^{th}$  user has power  $S_{nj}(t)$ .
- The system achieves an **instantaneous bit rate** for user  $n$  equal to :

$$R_n(t) = \Delta \log_2 \left( 1 + \frac{\sum_j L(X_n, Y_j) S_{nj}(t)}{\Delta \eta + \sum_k L(X_n, Y_k) \sum_{m \neq n} S_{mk}(t)} \right).$$

This condition is only a sufficient condition of achievability of the instantaneous bit rates.

## CDMA NETWORK IN MACRODIVERSITY

For CDMA networks :

$$\Delta \log_2(1 + SIR) \approx \frac{\Delta}{\log(2)} SIR$$

**Theorem 1.** *With the above approximation, if  $\pi$  is a policy, there exists a **finite power allocation** scheme such that :*

$$W_{t'} = W_t + \Sigma(t, t') - \sum_j \int_t^{t'} r_j \pi_j(s) ds,$$

$$r_j(x) = \frac{\Delta}{\log(2)} \frac{L(x, Y_j)}{I(x)}.$$

## STABILITY ISSUE

$$W_{t'} = W_t + \Sigma(t, t') - \sum_j \int_t^{t'} P_j(s) r_j \pi_j(s) ds.$$

A policy is **stable** if there exists a stationary workload satisfying the evolution equation.

The system is **stable** if there exists a stable policy.

For which class of arrival processes  $A$  is the system stable ?

## STABILITY REGION OF CELLULAR POLICIES

Let  $\{V_j\}_j$ , be a tessellation of  $\mathbb{R}^2$ , a **cellular policy with cells**  $\{V_j\}_j$  is a policy scheme satisfying for all  $j$  :

$$\forall m \in \mathcal{M}, \quad \pi_j(m)(\mathbb{R}^d \setminus V_j) = 0.$$

A cellular policy is **work-conserving** if  $m(V_j) > 0$  implies  $\pi_j(m)(V_j) = 1$ .

**Theorem 2.** *Any work-conserving cellular policy with bounded cells is stable if :*

$$\forall j, \quad \int_{V_j} \frac{\sigma(x)}{r_j(x)} \lambda(dx) < p_j.$$

*If there is a  $j$  such that :  $\int_{V_j} \frac{\sigma(x)}{r_j(x)} \lambda(dx) > p_j$  then any cellular policy with cells  $\{V_j\}_j$  is unstable.*

## STABILITY REGION : GENERAL CASE

$$\mathcal{F} = \left\{ f = (f_j)_j \text{ such that,} \right. \\ \left. \forall j, f_j \geq 0, \text{ measurable and } \lambda(dx)\text{-a.e. } \sum_j f_j(x) = 1 \right\}.$$

An  $f \in \mathcal{F}$  is a **partition of the plane**.

→ **Interpretation** : A user located at  $x$  has a part  $f_j(x)$  of its service provided by the  $j^{\text{th}}$  server station. For a cellular policy :  $f_j = \mathbb{I}(V_j)$ .

## STABILITY REGION

$$\mathcal{N}^s = \{A \in \mathcal{N} : \exists f \in \mathcal{F} \text{ such that } : \forall j, \int_{\mathbb{R}^2} \frac{\sigma(x) f_j(x)}{r_j(x)} \lambda(dx) < p_j\}.$$

$$\bar{\mathcal{N}}^s = \{A \in \mathcal{N} : \exists f \in \mathcal{F} \text{ such that } : \forall j, \int_{\mathbb{R}^2} \frac{\sigma(x) f_j(x)}{r_j(x)} \lambda(dx) \leq p_j\}.$$

**Theorem 3.** *For the spatial queuing model,*

- *if  $A \in \mathcal{N}^s$ , then there exists a stable policy,*
- *if there is a stable policy then  $A \in \bar{\mathcal{N}}^s$ .*

## OPTIMAL SPATIAL ALLOCATION

For simplicity, we suppose :  $\lambda(dx) = \lambda(x)dx$ . Stability relies on the value of :

$$\begin{aligned}\rho &= \inf_{f \in \mathcal{F}} \sup_{j \in \mathbb{N}} \frac{1}{p_j} \int_{\mathbb{R}^2} \frac{\sigma(x) f_j(x)}{r_j(x)} \lambda(x) dx \\ &= \inf_{f \in \mathcal{F}} \rho(f).\end{aligned}$$

If  $\rho < 1$ , the system is **stable**, if  $\rho > 1$ , the system is **unstable**.

## OPTIMAL SPATIAL ALLOCATION

Let

$$\mathcal{F}^* = \{f \in \mathcal{F} : \rho(f) = \rho\}.$$

**Proposition 1.** *There is an  $f \in \mathcal{F}^*$  such that :*

$$\forall j, \quad \frac{1}{p_j} \int_{\mathbb{R}^2} \frac{\sigma(x) f_j(x)}{r_j(x)} \lambda(x) dx = \rho.$$

*If  $\rho$  is finite and if there is a finite number of server stations, all  $f \in \mathcal{F}^*$  satisfy the above equation.*

→ optimal partitioning of space equalizes traffic load !

## OPTIMAL SPATIAL ALLOCATION

The extremal points of the convex set  $\mathcal{F}$  are the functions  $f_j = \mathbb{I}(V_j)$ , where  $\{V_j\}_j$  is a tessellation.

**Definition 2.** The processing rates are said to be *singular* if there exist  $j \neq k$ ,  $C > 0$  and a Borel set  $A$  of positive Lebesgue measure such that :

$$\forall x \in A, r_j(x) = Cr_k(x).$$

**Proposition 2.** If the processing rates are not singular, then *there is an  $f \in \mathcal{F}^*$  which is a tessellation*. If  $\rho$  is finite and if there is finite number of server stations, *all  $f \in \mathcal{F}^*$  are tessellations*.

→ For CDMA, Macrodiversity does not increase the Stability Region !

## REGULAR HEXAGONAL NETWORK

- Base Station on a regular hexagonal grid, with  $L$  km between two adjacent sites.
- Intensity of user arrival :  $\lambda$  users per square kilometers per time unit.
- A user requires to receive in mean  $\sigma$  bits of data from the network.
- The bandwidth is  $\Delta$  Hz on a CDMA Channel.
- The attenuation function from  $x$  to  $y$  is :  $K|x - y|^{-\alpha}, \alpha > 2$ .

## REGULAR HEXAGONAL NETWORK

$\lambda\sigma < \rho_c$  then the network is stable, if  $\lambda\sigma > \rho_c$  the network is unstable,

Macrodiversity and the usual CDMA hexagonal network have the same  $\rho_c$ .

$$\rho_c \approx \frac{2}{\log(2)\sqrt{3}} \frac{\Delta}{L^2} \left(1 + \frac{0.94}{\alpha - 2}\right)^{-1}.$$

$\lambda\sigma$  is the mean number of bits pumped per surface unit per time unit : this is the bit rate per surface unit of our network.

For  $\Delta = 1$  MHz,  $L = 1$  km,  $\alpha = 4$  we get : 1, 13 Mbits per square kilometer and per second.

## ERGODIC SITES OF SERVER STATIONS

- Server Station position is an **ergodic** spatial point process.
- $r_j(x - Y_j)$  is a mark of the server station point process.
- Intensity of user arrival :  $\lambda$  and mean service requirement :  $\sigma$  bits.

The system is stable if  $\lambda\sigma < \rho_c$  the maximal bit rate per surface unit of the system :

$$\rho_c^{-1} = \inf_{f \in \mathcal{F}} \sup_{j \in \mathbb{N}} \frac{1}{p_j} \int_{\mathbb{R}^2} \frac{f_j(x)}{r_j(x)} dx$$

→ The geometry of the network is contained in the processing rates. For TDMA and CDMA channels are function of  $\frac{l(|x - Y_j|)}{I(x)}$ .

## ERGODIC SITES OF SERVER STATIONS

From ergodicity :  $\rho_c$  is almost surely constant.

**Proposition 3.** *(Under a technical condition on  $r_j$  easily satisfied in practice) If the server stations position is a Poisson Point Process, then almost surely  $\rho_c = 0$ .*

## CONCLUSION

- The stability region defines the capacity of a network.
- For CDMA channel, under our hypothesis, macrodiversity does not increase the Stability Region on the Downlink. Similar negative results hold for macrodiversity in fixed rate CDMA channel (voice channel).
- On the contrary, macrodiversity has a huge impact on the Uplink.